• The scope of multithreading

LWPs, kernel and user threads

- kernel-level threads supported by the kernel
 - Solaris, Linux, Windows XP/2000
 - all scheduling, synchronization, thread structures maintained in kernel
 - could write apps using kernel threads, but would have to go to kernel for everything
- user-level threads supported by a user-level library
 - Pthreads, Java threads, Win32...
 - sched. & synch can often be done fully in user space;
 kernel doesn't need to know there are many user threads
 - problem with blocking on a system call

• LightWeight Processes - LWP

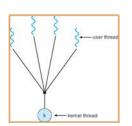
- these are "virtual CPUs", can be multiple per process
- the scheduler of a threads library schedules user-level threads to these virtual CPUs
- kernel threads implement LWPs => visible to the kernel, and can be scheduled
 - sometimes LWP & kernel threads used interchangeably, but there can be kernel threads without LWPs

Multithreading models

- There are three dominant models for thread libraries, each with its own trade-offs
 - many threads on one LWP (many-to-one)
 - one thread per LWP (one-to-one)
 - many threads on many LWPs (many-to-many)
- similar models can apply on scheduling kernel threads to real CPUs

Many to one

- In this model, the library maps all threads to a single lightweight process
- Advantages:
 - totally portable
 - easy to do with few systems dependencies
- Disadvantages:
 - cannot take advantage of parallelism
 - may have to block for synchronous I/O
 - there is a clever technique for avoiding it
- Mainly used in language systems, portable libraries



One to ne

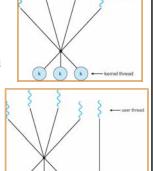
- In this model, the library maps each thread to a different lightweight process
- Advantages:
 - can exploit parallelism, blocking system calls



- thread creation involves LWP creation
- each thread takes up kernel resources
- limiting the number of total threads
- Used in LinuxThreads and other systems where LWP creation is not too expensive

Many to many

- In this model, the library has two kinds of threads: bound and unbound
 - bound threads are mapped each to a single lightweight process
 - unbound threads may be mapped to the same LWP
- Probably the best of both worlds
- Used in the Solaris implementation of Pthreads (and several other Unix implementations)



High Level Program Structure Ideas

- · Boss/workers model
- Pipeline model
- Up-calls
- Keeping shared information consistent using version stamps

Thread Design Patterns

Common ways of structuring programs using threads

- · Boss/workers model
 - boss gets assignments, dispatches tasks to workers
 - variants (thread pool, single thread per connection...)
 - Pipeline model
 - do some work, pass partial result to next thread
- Up-calls
 - fast control flow transfer for layered systems
- · Version stamps
 - technique for keeping information consistent

Boss/Workers

Boss: Worker:
forever {
get a request
switch(request)
case X: Fork (taskX)
case Y: Fork (taskY)
...
}

- · Advantage: simplicity
- Disadvantage: bound on number of workers, overheard of threads creation, contention if requests have interdependencies
- Variants: fixed thread pool (aka workpile, workqueue), producer/consumer relationship, workers determine what needs to be performed...

Pipeline

- Each thread completes portion of a task, and passes results
- like an assembly line or a processor pipeline
- Advantages: trivial synchronization, simplicity
- Disadvantages: limits degree of parallelism, throughput driven by slowest stage, handtuning needed

Up calls

- Layered applications, e.g. network protocol stacks have top-down and bottom-up flows
- Up-calls is a technique in which you structure layers so that they can expect calls from below
- Thread pool of specialized threads in each layer
 - essentially an up-call pipeline per connection
- Advantages: best when used with fast, synchronous control flow transfer mechanisms or program structuring tool
- Disadvantages: programming becomes more complicated, synchronization required for top-down

Version Stamps

- (not a programming structure idea but useful technique for any kind of distributed environment)
- maintain "version number" for shared data
 - keep local cached copy of data
 - check versions to determine if changed